## CSCB58: Computer Organization



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The content of this lecture is adapted from the lectures of
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## CSCB58 Week 11

## ||II <br> Question \#1

- How do you write an assembly language program that performs $\$ \mathrm{t} 0=\$ \mathrm{t} 1 \times \$ \mathrm{t} 2$ without using mult or multu?
- Coming up with a solution is easier if you ask yourself certain questions:
- How can multiplication be done using add?
- What if \$t2 stores a zero value?
- How do you make a loop happen?
- How do you make it stop looping?

What needs to be done at the beginning?

## Question \#1

- Assume, you'll have a list of available assembly language commands:

| Select |  | Input | Operation |  |
| :--- | :--- | :--- | :--- | :--- |
| $\mathrm{S}_{1}$ | $\mathrm{~S}_{0}$ | Y | $\mathrm{C}_{\mathrm{in}}=\mathbf{0}$ | $\mathrm{C}_{\text {in }}=\mathbf{1}$ |
| 0 | 0 | All 0 s | $\mathrm{G}=\mathrm{A}$ | $\mathrm{G}=\mathrm{A}+1$ |
| 0 | 1 | B | $\mathrm{G}=\mathrm{A}+\mathrm{B}$ | $\mathrm{G}=\mathrm{A}+\mathrm{B}+1$ |
| 1 | 0 | B | $\mathrm{G}=\mathrm{A}-\mathrm{B}-1$ | $\mathrm{G}=\mathrm{A}-\mathrm{B}$ |
| 1 | 1 | All 1s | $\mathrm{G}=\mathrm{A}-1$ | $\mathrm{G}=\mathrm{A}$ |

## Register table:

Register values: Processor role

- Register 0 (Szero): value 0
- Register 1 (Sat): reserved for the assembler
- Registers $2-3(\$ \mathrm{SVO}, \mathrm{Sv1})$ : return values
- Registers 4-7 (\$a0-Sa3): function arguments
- Registers 8-15, 24-25 (St0-St9): temporaries
- Registers $16-23$ (\$50-\$s7): saved temporaries
- Registers 28-31( $\mathrm{Sg} \mathrm{g}, \mathrm{S}, \mathrm{sp}, \mathrm{Sfp}, \mathrm{Sra})$



## Question \#1: The Math

- How can multiplication be done using add?

```
add $t0, $t0, $t1
(repeat this many times)
```

- What if \$t2 stores a zero value?

```
start: beq $t2, $zero, end
    ... # multiplication code here
end:
```


## Question \#1: The Loop

- How do you make the loop happen?

```
start:
end:
```

- How do you make it stop looping?

```
start: beq $t2, $zero, end
    addi $t2, $t2, -1
    j start
```

end:

## Question \#1: The combination

- What needs to be done at the beginning?
add \$t0, \$zero, \$zero
- Final solution:

```
add $t0, $zero, $zero
start: beq $t2, $zero, end
    add $t0, $t0, $t1
    addi $t2, $t2, -1
    j start
    end:
```


## Question \#2

- Final Exam, Winter 2012:

3. In the space below, write a short assembly language program that is a translation of the program on the right. You can assume that $i$ has been placed on the top of the stack, and that the return value should be placed on the stack as well before returning to the calling program. Make sure that you comment your code so that we understand what you're doing. (10 marks)
```
int sign (int i) {
    if (i > 0)
        return 1;
    else if (i < 0)
        return -1;
    else
            return 0;
```

- How would you convert this to assembly language?


## Question \#3: More Assembly

- Translate this C-style code into 4 lines of MIPS assembly code:

```
int t1= 10, t2=3;
int t3 = t1 + 2*t2
```

- Final solution:

```
li $t1, 10
li $t2, 3
sll $t2, $t2, 1
add $t3, $t2, $t1
```


## Question \#4: More Assembly

- Translate this C-style code into MIPS assembly code:

$$
\mathrm{A}[i]=\mathrm{A}[i / 2+1]+1 ;
$$

- Final solution:

| $l w$ | $\$ t 0, ~ 0(\$ g p)$ | $\#$ fetch $i$ |
| :--- | :--- | :--- |
| $l a$ | $\$ t 8, A$ | $\#$ fetch A |
| srl | $\$ t 1, \$ t 0,1$ | $\#$ i/2 |
| sll | $\$ t 1, \$ t 1,2$ | $\#$ turn i/2 into a byte offset (*4) |
| add | $\$ t 1, \$ t 8, \$ t 1$ | $\#$ \&A[i/2] |
| $l w$ | $\$ t 1,4(\$ t 1)$ | $\#$ fetch A[i/2 +1$]$ |
| addi | $\$ t 1, \$ t 1,1$ | $\#$ A[i/2 +1$]+1$ |
| sll | $\$ t 2, \$ t 0,2$ | $\#$ turn i into a byte offset |
| add | $\$ t 2, \$ t 8, \$ t 2$ | $\#$ \&A[i] |
| sw | $\$ t 1,0(\$ t 2)$ | $\# A[i]=\ldots$ |

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