CSCB58: Computer Organization



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University of Toronto Fall 2020



The content of this lecture is adapted from the lectures of Larry Zheng and Steve Engels

CSCB58 Week 11

Question #1

- How do you write an assembly language program that performs \$t0 = \$t1 x \$t2 without using mult or multu?
- Coming up with a solution is easier if you ask yourself certain questions:
 - How can multiplication be done using add?
 - What if \$t2 stores a zero value?
 - How do you make a loop happen?
 - How do you make it stop looping?
 - What needs to be done at the beginning?

Question #1

 Assume, you'll have a list of available assembly language commands:

Reference Information

ALU arithmetic input table:

Select		Input	Operation	
S ₁	S ₀	Y	Ç _{in} =0	Ç _{in} =1
0	0	All Os	G=A	G=A+1
0	1	В	G=A+B	G=A+B+1
1	0	В	G=A-B-1	G=A-B
1	1	All 1s	G=A-1	G=A

Register table:

Register values : Processor role
 Register 0 (\$zero): value 0.

- Register 1 (\$at): reserved for the assembler.
- Registers 2-3 (\$v0, \$v1): return values
- Registers 4-7 (\$a0-\$a3): function arguments
- Registers 8-15, 24-25 (\$t0-\$t9): temporaries
- Registers 16-23 (\$s0-\$s7): saved temporaries

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Registers 28-31 (\$gp, \$sp, \$fp, \$ra)

mstruction	Op/rune	Syntax
add	100000	\$d, \$s, \$t
addu	100001	\$d, \$s, \$t
addi	001000	\$t, \$s, į
addiu	001001	\$t, \$s, j
div	011010	\$s, \$t
divu	011011	\$s, \$t
mult	011000	\$s, \$t
multu	011001	\$s, \$t
sub	100010	\$d, \$s, \$t
subu	100011	\$d, \$s, \$t
and	100100	\$d, \$s, \$t
andi	001100	st, ss, į
nor	100111	\$d, \$s, \$t
or	100101	\$d, \$s, \$t
ori	001101	\$t, \$s, į
XOL	100110	\$d, \$s, \$t
xori	001110	\$ t, \$ s, į
311	000000	\$d, \$t, a
ally	000100	\$d, \$t, \$s
ara	000011	\$d, \$t, a
sray.	000111	\$d, \$t, \$s
srl	000010	\$d, \$t, a
arly	000110	\$d, \$t, \$s
beg	000100	\$s, \$t, label
batz.	000111	\$s, label
blez	000110	\$s, label
bne.	000101	\$s, \$t, label
j	000010	label
jal	000011	label
jalr	001001	\$ 5
jr	001000	\$5
lb	100000	\$t, į (\$s)
lbu	100100	\$t, į (\$s)
lh	100001	\$t, į (\$s)
lhu	100101	\$t, į (\$s)
lw	100011	\$t, į (\$s)
3b	101000	\$t, į (\$s)
sh	101001	\$t, į (\$s)
<u>314</u>	101011	\$t, į (\$s)
trap	011010	į
mflo	010010	\$d

Instruction table:

Instruction On/Fune Suntay

Student Number: _____

Question #1: The Math

How can multiplication be done using add?

add \$t0, \$t0, \$t1

(repeat this many times)

What if \$t2 stores a zero value?

start:	beq \$t2, \$zero, end
	• • •
	# multiplication code here
	• • •
end:	•••

Question #1: The Loop

How do you make the loop happen?



How do you make it stop looping?



Question #1: The combination

What needs to be done at the beginning?

add \$t0, \$zero, \$zero

Final solution:

	add Stu, Szero, Szero
start:	beq \$t2, \$zero, end
	add \$t0, \$t0, \$t1
	addi \$t2, \$t2, -1
	j start
end:	• • •

Question #2

Final Exam, Winter 2012:

3. In the space below, write a short assembly language program that is a translation of the program on the right. You can assume that i has been placed on the top of the stack, and that the return value should be placed on the stack as well before returning to the calling program. Make sure that you comment your code so that we understand what you're doing. **(10 marks)**

int sign (int i) {
 if (i > 0)
 return 1;
 else if (i < 0)
 return -1;
 else
 return 0;</pre>

How would you convert this to assembly language?

Question #3: More Assembly

Translate this C-style code into 4 lines of MIPS assembly code:

int t1= 10, t2=3; int t3 = t1 + 2*t2

Final solution:

li \$t1, 10 li \$t2, 3 sll \$t2, \$t2, 1 add \$t3, \$t2, \$t1

Question #4: More Assembly

Translate this C-style code into MIPS assembly code:

```
A[i] = A[i/2 + 1] + 1;
```

Final solution:

```
lw $t0, 0($gp)  # fetch i
la $t8, A  # fetch A
srl $t1, $t0, 1  # i/2
sll $t1, $t1, 2  # turn i/2 into a byte offset (*4)
add $t1, $t8, $t1  # &A[i/2]
lw $t1, 4($t1)  # fetch A[i/2 + 1]
addi $t1, $t1, 1  # A[i/2 + 1] + 1
sll $t2, $t0, 2  # turn i into a byte offset
add $t2, $t8, $t2  # &A[i]
sw $t1, 0($t2)  # A[i] = ...
```

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